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AN AUTOMATIC INTEGRATOR FOR ORDINARY
DIFFERENTIAL EQUATIONS FOR ILLIAC IV

by

Thomas McCarthy

June 1, 1969

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DEPARTMENT OF COMPUTER SCIENCE
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Report No. 336

AN AUTOMATIC INTEGRATOR FOR
ORDINARY DIFFERENTIAL EQUATIONS FOR ILLIAC IV

by

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June 1, 1969

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ABSTRACT

This report describes the design and implementation of an automatic integrator for systems of ordinary differential equations for use on the ILLIAC IV computer system. This integrator accepts a simple statement of the differential equations and their initial values and produces an ILLIAC IV Assembly Language code which can be used to solve them. The intent is to produce a code which will be efficient on a large parallel computer.

ACKNOWLEDGMENT

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1. INTRODUCTION

Standard methods of integrating ordinary differential equations consist essentially of incrementing the independent variable by some small amount and then, by using past values and derivatives of the dependent variable, predicting the new value of that dependent variable. Thus the integration proceeds from the given initial value of the dependent variable to some user specified stopping point, building on itself as it goes. The same method is applied to systems of ordinary differential equations; but, of course, one cannot integrate one equation completely and then the next, since in general, the derivative of each dependent variable will depend on the other dependent variables of the system. Each equation must be integrated at a given value of the independent variable before going to the next step.

For this reason, to perform an integration of a system on a large parallel computer like ILLIAC IV, it is desirable to perform the integration on all of the equations simultaneously. Difficulties arise when the equations are not identical in form; and, of course, in general they are not. Integration methods consist of one or more function evaluations; that is, substituting the present value of the dependent variable into the equation to find its derivative at the given point. This is the part that is most difficult to do simultaneously with dissimilar equations, and it is the part which usually accounts for much of the computation time, even on a serial machine.

The following example will illustrate some of the problems to be solved:

$$\begin{aligned} y_1' &= a y_1 y_2 + b y_1 + c y_1 y_1 / y_2 \\ y_2' &= d y_2 + e y_1 y_2 + f y_1 y_2 \end{aligned}$$

The logical way to proceed would seem to be to try to minimize the number of nonparallel arithmetic operations which must be done. For an example of how this might be done, consider a two processing element machine with each processor working on a single equation. Assume that, by allocating storage properly and by using some indexing (to be explained later), it is possible to access constants and different dependent variables simultaneously. The operation sequence would then be as follows (parentheses indicate memory access, I indicates processor idle, arithmetic operations indicated by their symbols).

```
P1 (a) x x (b) x I + (c) x x ÷ +
P2 (d) x I (e) x x + (f) x x I +
```

If the terms can be rearranged so that the first term of the first equation and the second term of the second equation can be done at the same time, some more savings can be gained.

```
P1 (a) x x (b) x + (c) x x ÷ +
P2 (e) x x (d) x + (f) x x I +
```

Thus more operations are done in parallel and the processors are idle fewer times. With larger systems of equations there are more terms that can be done in parallel and more comparisons to be made. With rearrangement, unmatched terms could also be done partly in parallel. With the larger number of terms in big systems, it is reasonable to expect that the relative savings will be greater than those in this example.

Another problem would arise if one of the equations had significantly more terms than the other. Some method of sharing the work between processors should be used. As the system of equations gets larger in number of terms and number of equations, all of these problems become more and more complex.

The translator-integrator program described in this thesis tries to do these types of optimizations as well as compiling the Assembly code to evaluate the input equations. The translator-integrator (which will be referred to as the integrator in the following pages - the section designated should be clear from the context) has associated with it some skeleton forms of ILLIAC IV Assembly Language programs to solve ordinary differential equations by different numerical methods. These, along with the compiled code and generated storage pseudo-ops allows the program to produce a completed Assembly code to solve the equations.

The translator-integrator system itself is an Algol program which would be executed on the Burroughs 6500 in the ILLIAC IV system. A schematic of the integrator is shown in Figure 1.

Section 2 describes the integrator from the user's point of view. Section 3 is a discussion of the implementation techniques and numerical methods used. Section 4 discusses the capabilities of this system and the further work which might be done.

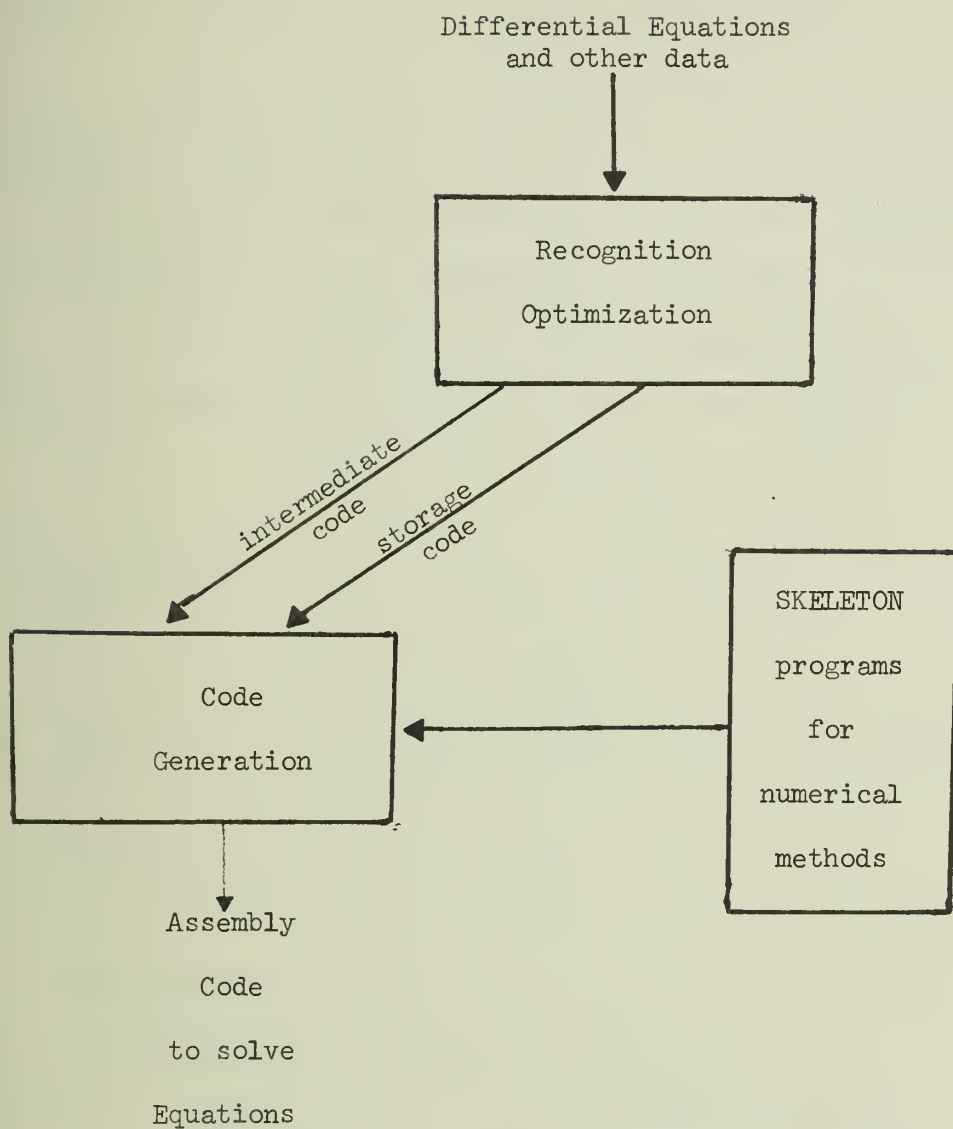


Figure 1. Schematic Diagram of Integrator System

2. USE OF THE INTEGRATOR SYSTEM

The primary object of any automatic integration system is to provide the user with a device to solve his equations with a minimum of effort and analysis. From the user's point of view, the integrator input is as defined in Backus Normal Form (BNF) in Figure 2. Most of this should be self-explanatory. Dependent variables are expressed as $Y(N)$; their derivatives as $Z(N)$, and the independent variable as T . Each differential equation is written as

$$Z(N) = f_n(Y(1), Y(2), \dots, T)$$

The functions can be almost any Algol expression (an obvious exception is an if clause) as defined by the BNF and is written in B6500 Algol notation. It should be noted that constants whose values can be assigned later are allowed in these expressions. The initial values of the dependent variables and the starting and final values of the independent variable are then given. These are real numbers expressed in Algol notation. All numbers except indexes are treated as 64-bit floating point numbers by the integrator. It is not necessary to terminate whole numbers with a decimal point to insure that they are floating point. T_0 and T_F are the starting and ending values of T , the independent variable.

The user can also either specify a stepsize to be used or specify a desired maximum error which the integrator will try to stay

```

<program> ::= <differential equations> <initial values> <constant declarations>
           <other statements> END

<differential equations> ::= Z( <integer> ) = <expression> ; |
           <differential equations> Z( <integer> ) = <expression> ;

<initial values> ::= ZO( <integer> ) = <real number> ; |
           <initial values> ZO( <integer> ) = <real number> ;

<constant declarations> ::= <identifier> = <real number> ; |
           <constant declarations> <identifier> = <real number> ;

<other statements> ::= <statement> ; | <other statements> <statement> ;

<statement> ::= TO = <real number> | TF = <real number> | STEP = <real number> |
           ERROR = <real number>

<expression> ::= <term> | <add operator> <term> | <expression> <add operator>
           <term>

<add operator> ::= + | -

<term> ::= <factor> | <term> <multiply operator> <factor>

<multiply operator> ::= x | /

<factor> ::= <primary> | <factor> * <primary>

<primary> ::= <real number> | <identifier> | Y( <integer> ) | T |
           <funid> ( <expression> ) | ( <expression> )

<funid> ::= SIN | COS | EXP | LOG

```

Figure 2. Input Language Description in BNF

within. Also any constants which were used in the equations should be defined. The program input is terminated with an END.

In general the input is free form with semicolons separating statements. In fact, a missing semicolon will not always cause an error, but it is suggested that the user try to make the input no more difficult to recognize than necessary. Identifiers are limited to six characters; longer ones are truncated on the right. Numbers are limited to eighteen characters including sign and exponent. All blanks are ignored.

From this minimal amount of information an Assembly Language Code will be produced and placed as a file on disk. If some information, such as an initial value or constant declaration is omitted, a code will still be produced with zeros in the appropriate locations. In Figure 3, a sample program for the example discussed in the introduction is shown as it would be inputted to the translator. A larger example, complete with output, is given in the Appendix.

The integrator also has some control words which can be used if desired. These are entered at the beginning of the input. The string of control words separated from each other by at least one space must be preceded by a '#' and a space. The control words are as follows:

TEST - Do not generate Assembly Language Code. This control word is used when the user is only interested in knowing what efficiency can be achieved with his equations.


```

      # RUNGE  ERRORCHECK      ;

Z(1)=  AxY(1)xY(2) + BxY(1)  +  CxY(1)xY(1)/Y(2)  ;

Z(2)=  DxY(2)+ExY(1)xY(2) + FxY(1)xY(2)  ;

ZO(1)=  0  ;

ZO(2)=  2.3  ;

A = 3.456  ;

B = 0  ;

C = 1  ;

D = 3.45677483645@-09  ;

E = 4.9  ;  F = .33  ;

TO = 0.0  ;  TF = 1.0  ;

ERROR = .0001  ;

END

```

Figure 3. Sample Input

- RUNGE - Use a Fourth-Order Runge Kutta method to integrate the equations. This is the default option.

- CORPRE - Use a Corrector-Predictor method to integrate the equations. This is a more complicated method but could be used on more difficult equations.

- ERRORCHECK - Estimate truncation error and adjust stepsize accordingly. This control word should be used when an error bound is to be put on the results as previously described.

- CONSTMAP - This prints out a chart showing the location of constants in PE memory so that they can be changed without reexecuting the integrator.

- FUNMODE - If the terms in the differential equations are predominantly made up of parenthesized expressions or function calls (i.e., SIN, COS, etc.) or both, then this option should be turned on to improve the efficiency of the code generated.

- PATTERN - The integrator makes some of its decisions based on the supposition that the terms in the first few equations are representative types for the whole system. If this is not so for a particular system, the efficiency may be improved by using this control word. When this option is used, the first equation should contain terms typical of the system and in about the same ratio to the other terms in the equation as they occur in the system. No code will be generated for this equation, but it will aid the integrator in allocating the work to be done.

The Assembly Language Program which is produced also has some comments in it so that with the information provided by the CONSTMAP option of the integrator, the user should be able to make small changes, such as the values of constants or initial values, without rerunning the program if desired. An alternative to this would be to

read in data at execution time. When the input conventions for ILLIAC IV are defined, the integrator can easily be modified to load the data into the correct place.

3. IMPLEMENTATION

3.1 An Overview of the System

3.1.1 The Recognizer

The recognizing part of the integrator program was written using the ILLIAC IV Translator Writing System (TWS). The input to one of the TWS programs is a Backus Normal Form (BNF) description of the integrator input language syntax. This program, which runs on the Burroughs 5500, generates an Algol scanner and parser which can be used to recognize programs written in the language defined by the input BNF. An action number can be associated with each production to execute a designated semantic routine. A TWS program is also provided to concatenate these semantic routines to the parser. While this program also provides several built-in routines designed to help write compilers, few of these were actually used due to the specialized nature of the integrator language. All of the analysis, reduction and code generation for the integrator is done within the semantic routines.

The actual input to the TWS programs is somewhat different than that given in Figure 2. Action numbers are required, and the BNF description is slightly different to aid the syntactic analysis.

3.1.2 The Semantic Routines

Before the parsing of the input begins, a short, fast prepass is made in the initial semantic routine to count the number of equations

and number of terms (by counting arithmetic operators). Since the input to the integrator may be quite large, it is necessary to do optimization and the related compilation and storage allocation as the program proceeds, so as to limit the amount of B6500 storage needed. The information gleaned from this prepass is needed to do this dynamic analysis.

As the input differential equations are parsed, an intermediate code for evaluating them is generated. The optimization is then performed on a term by term basis on this intermediate code. Storage assignment begins at the first processing element and then proceeds according to the types of terms and equations involved in the problem. Other types of statements such as initial values and constant declarations are recognized, and the information is saved. When the entire input has been recognized, all the necessary information for Assembly code generation has been formed. At this point the appropriate skeleton ILLIAC IV integration program is selected, and the storage declarations and the code for evaluating the equations are inserted in the proper places to form a completed Assembly Language Code. This is left as a disk file which can be assembled to machine code.

3.2 Methods for Improving Parallelism

3.2.1 Possible Instructions and Operands

The instructions used to evaluate an equation are a small and standard set; for example, load register, add to register, multiply,

store, etc. Thus the operands that will be used also constitute a fairly limited set. In general, they can be divided into the following classifications:

1. **CONSTANTS** - These values do not change during the integration. They may have been declared either as literals or by constant declaration statements.
2. **DEPENDENT VARIABLE** - These have an index value associated with them declared in input language as Y (<index>).
3. **FUNCTION CALLS** - These include any of the explicitly allowed functions such as SIN and COS.
4. **INDEPENDENT VARIABLE** - This is designated by T in the input language.
5. **REGISTER NAME** - These occur in register to register instructions used for temporary storage of data.
6. **NO OPERAND** - This is for such instructions as change the sign of the A register (CHSA).

3.2.2 Doing Similar Instructions in Parallel

First consider the same instruction operating on different operands which fall into the same operand class. Methods must be developed for doing these in parallel.

When the operand is a constant, as many consecutive rows of PE memory as needed are blocked off. When the Assembly code is generated, instructions which access constants are given operand fields such that they point to consecutive rows in this array. Each processor can have a different value in its memory position in the constant row.

It is a simple matter for the integrator to take the values defined as literals or by constant declarations and insert them in the correct row and column so they are accessed by the correct instruction.

The values of the dependent variables are not stored in the same way as constants since they must be updated constantly throughout the execution of the integration to be performed. The values of the independent variables that are needed are kept in each PE in a column. The way to access different Y's in parallel is to associate a tag row with each access in a way similar to that used for constants. In this case the tag row is loaded into the X register when needed and a load modified by index is made from the independent variable storage area. While this costs an instruction to load the X register, it really only costs $1/m$, where m is the number of equations, since the indexes for all of the variables are loaded at once. Also in a serial machine these variables would often be declared as an indexed array.

At the present stage of design, it is not possible to execute different function calls in parallel, since a function call is in reality a reasonably large set of instructions to be inserted in this position, and the processing elements must execute the same instructions. Since most functions form some sort of polynomial approximation, it would be a reasonable extension of the system to have the integrator provide its own library functions which have been written so they are as parallel as possible and thus gain some speed.

Since there is only one independent variable, instructions with it as an operand can always be performed in parallel. The value is

kept in all PE's in the same memory row and updated when appropriate. No operand instructions can clearly always be done in parallel.

Instructions with different register names as operands clearly cannot be done in parallel. Since these only occur in temporary storage manipulation according to a well-defined algorithm, similar terms will probably access the same registers anyway. Also, due to the method used for sharing terms described in a later section, little temporary storage will be used.

3.2.3 Term by Term Analysis

In order to compare equations to see what can be done in parallel, it is necessary to break them down in to more convenient and hopefully more similar parts. The unit selected was that of a term since these can be shared to other processors and then recombined with just an addition operation. The expressions representing the differential equations are only broken down into the first level of terms; that is, terms inside parentheses are not separated. Some examples are listed below. The expressions on the left are broken down into the terms (separated by commas) on the right. The expression can be reformed by just adding the terms on the right.

$$A + B + C \longrightarrow A, B, C$$

$$A + B - C \longrightarrow A, B, -C$$

$$A + (B-C) \longrightarrow A, (B-C)$$

$$AxY(1) + BxY(2) \times Y(3) + \text{SIN}(Y(4)) \longrightarrow \\ AxY(1), BxY(2) \times Y(3), \text{SIN}(Y(4))$$

Of course, one can certainly write separate terms according to this rule which differ tremendously, but it is reasonable to expect that in a large system of equations many terms will be similar. Further discussion of the integrator should make obvious the convenience of this division. Thus if all of the instructions in two terms are the same and the operands fall into the same classes as defined by previous remarks, then these terms can be done in separate PE's simultaneously.

3.2.4 Inserting Dummy Instructions

In actuality it cannot be assumed that all terms will be perfect matches with other terms. Consider the following terms:

$$Ax \ Y(1) \times Y(3)$$

$$Bx \ Y(3)$$

They differ only in one multiply; and if no other better pairing is available, this can be treated as $Bx \ Y(3) \times 1$. In actuality it is treated as $Bx \ Y(3) \times Y(0)$ where $Y(0)$ always contains a one. Thus these terms can be paired. Other dummy instructions can be inserted as follows:

$$\left. \begin{array}{l} Y(1) \times B \\ Y(2) \end{array} \right\} \rightarrow \begin{array}{l} Y(1) \times B \\ Y(2) \times 1 \end{array}$$

$$\left. \begin{array}{l} \text{SIN}(A + Y(1)) \\ \text{SIN}(B) \end{array} \right\} \rightarrow \begin{array}{l} \text{SIN}(A + Y(1)) \\ \text{SIN}(B + 0) \end{array}$$

$$\left. \begin{array}{l} - (A + B) \\ - C \end{array} \right\} \longrightarrow \begin{array}{l} - (A + B) \\ - (C + 0) \end{array}$$

Inserting dummy operands for instructions involving other than constants and dependent variables does not appear worthwhile since extra instructions would be needed. Since a function call usually involves several instructions, it is not advisable to have a dummy call unless no better match can be found.

3.2.5 Sharing Terms

As previously mentioned, one of the important problems is that the equations may vary greatly in the number of terms in each. Some method is needed to share the work more or less equally between processors. This is the place where the information gained in the prepass becomes necessary. A limit is set on the number of terms to be evaluated in each PE according to the following equation:

$$\begin{array}{l} \text{maximum number} \\ \text{of terms / PE} \end{array} = \left[\frac{\text{total number of terms}}{\text{total number of PE's}} \right]$$

The square brackets indicate the smallest integer greater than the enclosed value. This limit serves to distribute the terms equally across the PE's. When the optimization has assigned the maximum number of terms to a PE, it then skips to the next. This requires that a tag be associated with each term indicating which equation it belongs to.

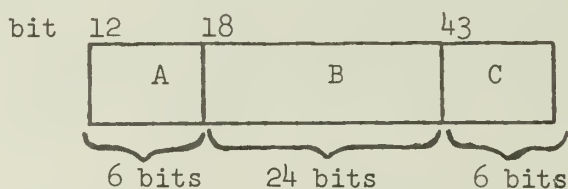
After a term has been evaluated, it is added to a scratch memory location in the same PE. This location depends on the tag so that at the end of the evaluation, there may be several separate partial sums in a particular PE. After all partial sums are formed, routes and adds are performed which combine partial sums and return them to the correct PE's. Also, if the number of PE's is much greater than (i.e., more than twice) the number of equations, the integration procedures other than function evaluation are assigned to every other one or two PE's sequentially so that the amount of routing necessary will be reduced. Assigning the PE's sequentially assures that the routes will be small since the terms are also assigned to the PE's more or less sequentially within the range of the optimization routine.

3.3 The Optimization Routine

The optimization routine is the procedure which is called when a term has been recognized. Its purpose is to determine what term this just recognized term can be done in parallel with. Flowcharts are given in Figures 4-6. There are three basic stacks. COMPSTK is the stack containing the intermediate code for the term just recognized. CODESTK contains the intermediate code for the terms which are to be shared over the processors as previously discussed. EXCODESTK contains the code for the terms which cannot be matched with those in CODESTK. In the normal state of operation, terms which contain function calls and parenthesized expressions are put in this stack since they are expected

to be the exception and not the rule. This can be changed by the user if desired with the control word FUNMODE. Optimization does occur within EXCODESTK as well as CODESTK.

Basically the optimization routine executes a series of calls to a comparison routine with different arguments. The arguments tell this comparison routine which stacks to compare and how many dummy instructions it can insert in each in order to make terms match. The flowcharts show the sequence of these calls. The purpose of the particular order is to find the match which requires the least overhead first. When the maximum number of terms allowed in CODESTK is reached, the data for storage generation is extracted for that PE, and the program proceeds to the next. CODESTK and EXCODESTK are overwritten for each PE since the basic instruction and operand will be the same if a comparison is achieved. Only the storage information contained in the intermediate code is different. The pattern for the intermediate code given below illustrates this.



A - code number indicating the instruction
(LDA, ADRN, . . .)

B - storage information

C - type of operand

If C is	→	then B contains
No operand	→	0
Constant	→	symbol table pointer
Dependent Variable	→	index value
Independent Variable	→	0
Function Call	→	Name of function
Register to Register	→	Source register name
Arithmetic computation temporary storage	→	index within storage
Exponentiation temporary storage	→	index within storage

The code remaining when the final equation has been optimized is used to form the ILLIAC IV code itself.

An example of the previously discussed case follows:

EQUATIONS:

$$Z(1) = AxY(1) \times Y(2) + BxY(1) + CxY(1) \times Y(2) / Y(2)$$

$$Z(2) = DxY(2) + ExY(2) \times Y(2) + FxY(1) \times Y(2)$$

DETERMINED IN PASS 1:

MAXIMUM TERMS PER PE = 3

NUMBER OF EQUATIONS = 2

NOTE: ASSUMING A 2 PE SYSTEM

PASS 2:

1. FIRST THREE TERMS ARE READ IN TO FILL UP CODESTK.
STORAGE ALLOCATION FOR THE FIRST PE IS DONE.

$$\left. \begin{array}{l} / Y(2) \\ x Y(1) \\ x Y(1) \\ C \end{array} \right\} \text{TERM \#3}$$

$$\left. \begin{array}{l} x Y(1) \\ B \end{array} \right\} \text{TERM \#2}$$

$$\left. \begin{array}{l} x Y(2) \\ x Y(1) \\ A \end{array} \right\} \text{TERM \#1}$$

CODESTK

2. READ NEXT TERM IN INPUT AND PUT IN COMPSTK.

$$\begin{array}{l} x Y(2) \\ D \end{array}$$

COMPSTK

3. COMPARE COMPSTK TO CODESTK WITH NO DUMMY INSERTIONS ALLOWED.
4. FIND A MATCH WITH TERM 2.

5. COMPSTK IS OVERWRITTEN INTO CODESTK. FLAGS ARE SET INDICATING THAT THIS TERM IS TAKEN.

6. NEXT TERM IS READ INTO COMPSTK.

x Y(2)
x Y(1)
E

COMPSTK

7. COMPARE COMPSTK TO CODESTK WITH NO DUMMY INSERTIONS ALLOWED.

8. FIND A MATCH WITH TERM #1.

9. OVERWRITE AS IN STEP 5.

10. READ IN NEXT TERM.

x Y(2)
x Y(1)
F

COMPSTK

11. COMPARE WITH REMAINING TERM IN CODESTK - NO INSERTIONS.

12. FAILURE - TRY TO COMPARE ALLOWING ONE INSERTION

13. SUCCESS

```

/ Y(0)
x Y(2)
x Y(1)
F

```

Note: Y(0) always contains
a 1.

COMPSTK

14. OVERWRITE AS IN STEP 5.

```

/ Y(0)
x Y(2)
x Y(1)
F

```

```

x Y(2)
D

```

```

x Y(2)
x Y(1)
E

```

CODESTK

15. USE CODESTK TO ALLOCATE STORAGE FOR THIS PE.

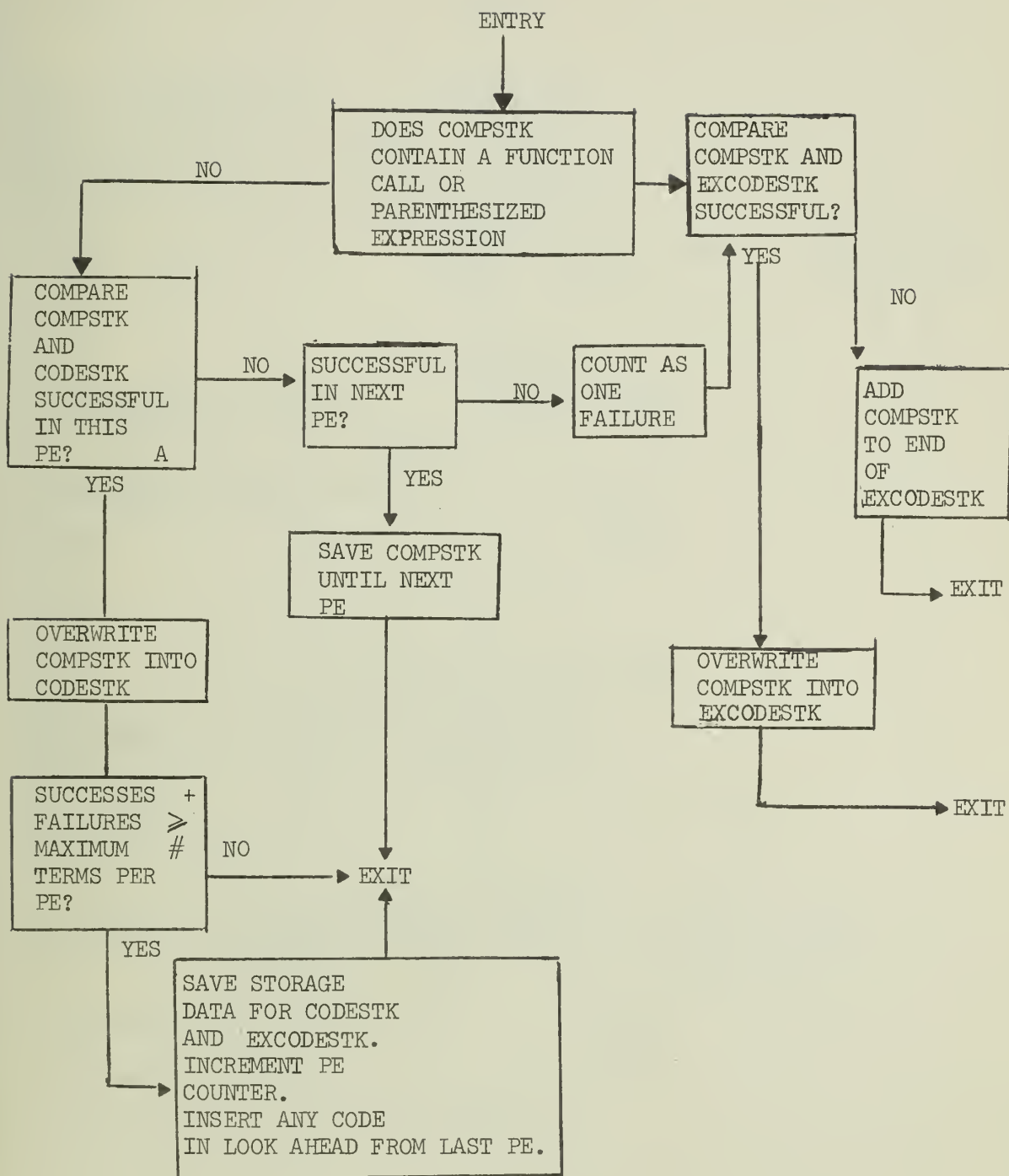


Figure 4. Optimization Routine

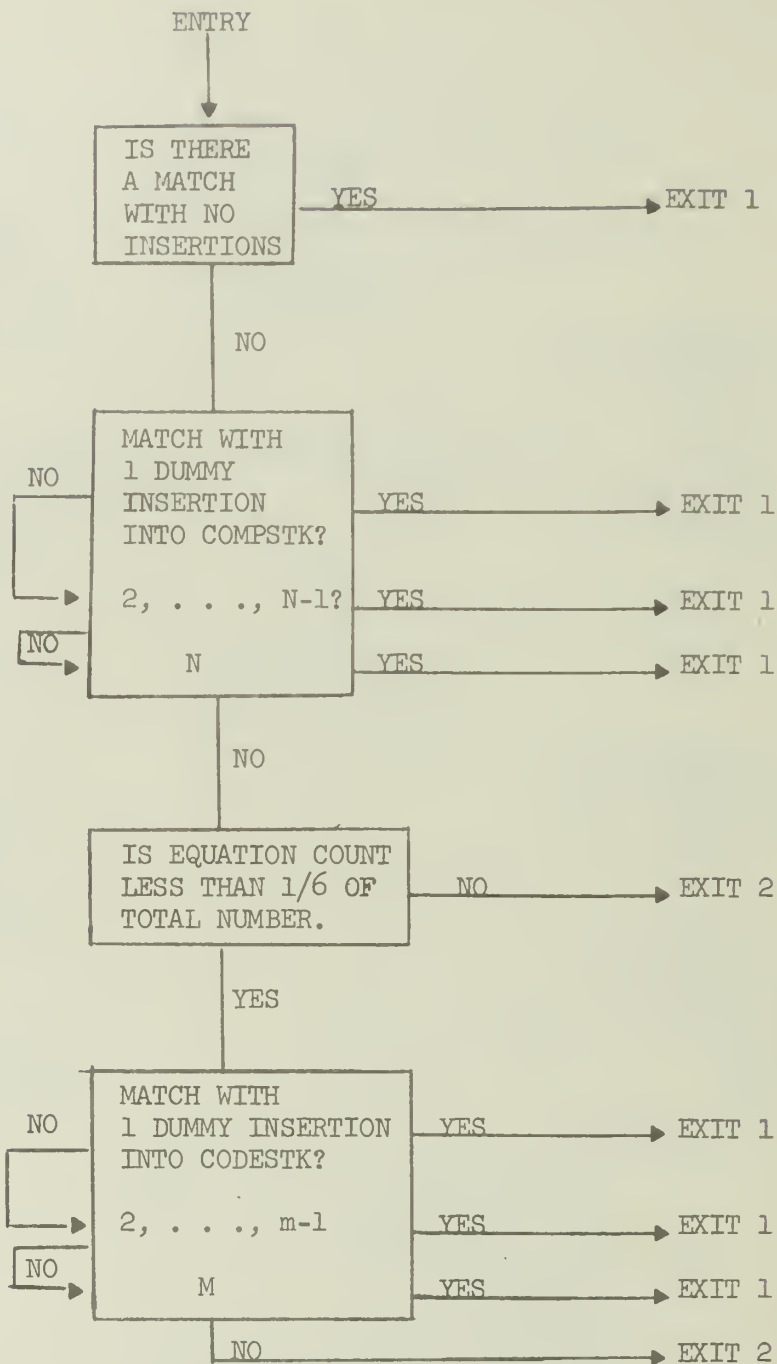


Figure 5. Comparing Compstk and Codestk BLOCK A in PREVIOUS FLOWCHART

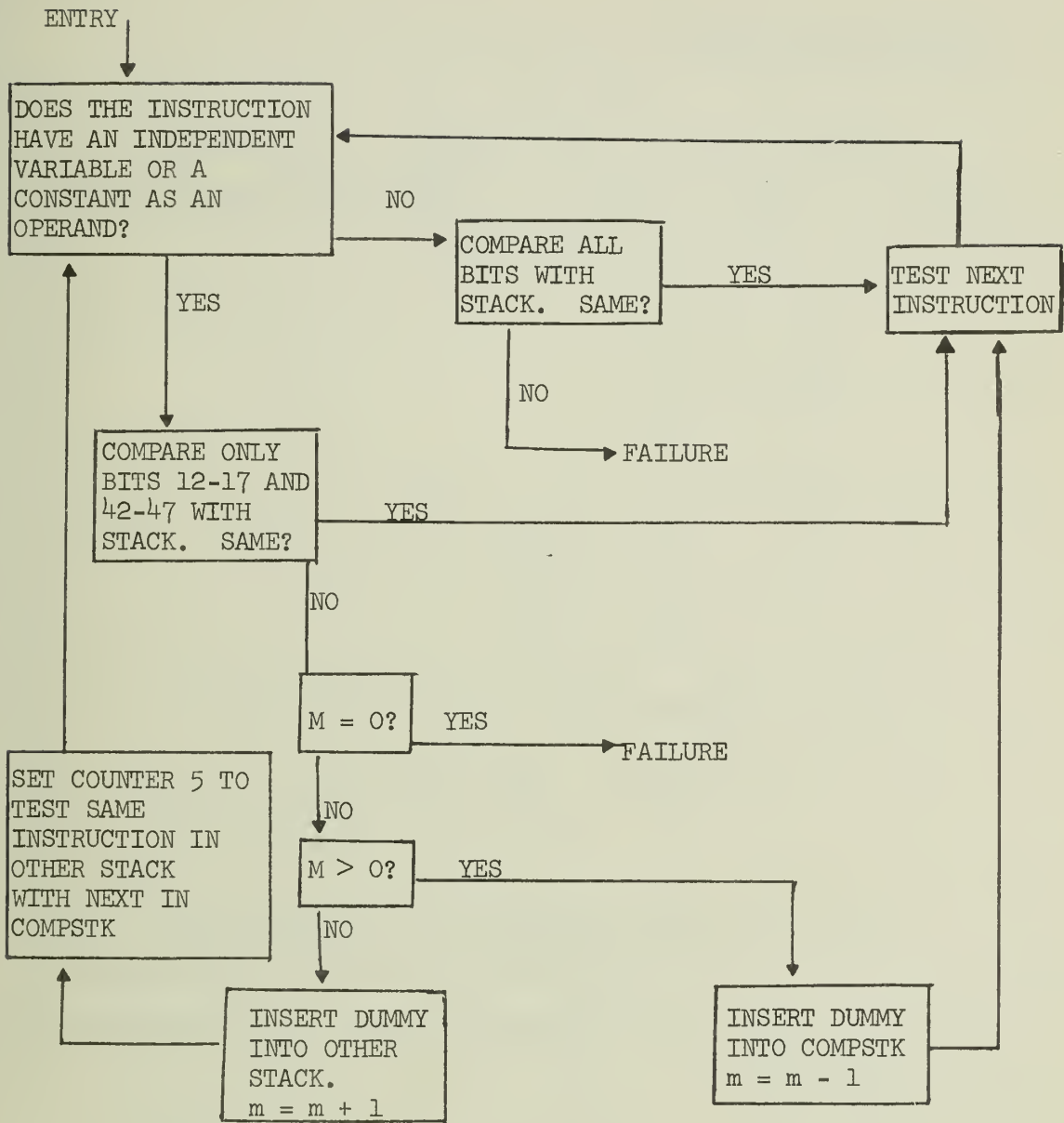


Figure 6. Routine to Compare COMPSTK with Another Stack. M is the number of insertions of dummy instructions allowed.

3.4 Methods of Integration

At present there are two integration methods available for the integrator. These are Fourth Order Runge Kutta and a corrector-predictor method investigated by Nordsieck.

The standard Fourth Order Runge Kutta method is used. This is as in the equations below.

$$K_1 = F(Y_n, T)$$

$$K_2 = F(Y_n + \frac{1}{2} HK_1, T + H/2)$$

$$K_3 = F(Y_n + \frac{1}{2} HK_2, T + H/2)$$

$$K_4 = F(Y_n + HK_3, T + H)$$

$$Y_{n+1} = Y_n + \frac{1}{6} H (K_1 + 2K_2 + 2K_3 + K_4)$$

If automatic stepsize adjustment is called for, spare PE's, which are not used for the integration, are used to do the integration at twice the stepsize. The truncation error can then be estimated from the following formula.

$$E_t = Kh^5 = \frac{32}{31} \left[y_m^{(h/2)} - y_m^{(h)} \right]$$

Stepsizes are always changed by a factor of two and certain limitations are put on the changing frequency and minimum size of the step.

The corrector-predictor method is a multistep method and thus requires another method to initialize it. The midpoint rule is used in the integrator. When the starting values y_{n-1} and y_{n-2} are found, they are premultiplied by a matrix as follows:

$$\underline{a}_n = \begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 3/4 & -1 & 1/4 \\ 0 & 1/6 & -1/3 & 1/6 \end{bmatrix} \begin{bmatrix} y_n \\ hf(y_n) \\ hf(y_{n-1}) \\ hf(y_{n-2}) \end{bmatrix}$$

It can be shown that the resulting \underline{a}_n is equivalent to

$$\left[y_n, hy_n', h^2 y_n''/2, h^3 y_n'''/6 \right]^T$$

The formulas for the integration method are then:

$$\underline{a}_{n+1} = \begin{bmatrix} 1 & 1 & 1 & 1 \\ 0 & 1 & 2 & 3 \\ 0 & 0 & 1 & 3 \\ 0 & 0 & 0 & 1 \end{bmatrix} \underline{a}_n + \begin{bmatrix} -3/8 \\ -1 \\ -3/4 \\ -1/6 \end{bmatrix}$$

where $F = h y_n' - h f(y_n)$. There are several reasons for using this method. For reasons which are beyond the scope of this explanation,

the method minimizes roundoff error. Also, due to the Pascal triangle form of the matrix, the matrix multiplication can be reduced to a series of a few additions thus speeding up the method. The biggest advantage is that changing the stepsize does not require restarting procedures.

The vector \underline{a}_n simply is multiplied by

$$\begin{bmatrix} 1 & 0 & 0 & 0 \\ 0 & \alpha & 0 & 0 \\ 0 & 0 & \alpha^2 & 0 \\ 0 & 0 & 0 & \alpha^3 \end{bmatrix} \quad \text{where } \alpha = \frac{h_{\text{old}}}{h_{\text{new}}}$$

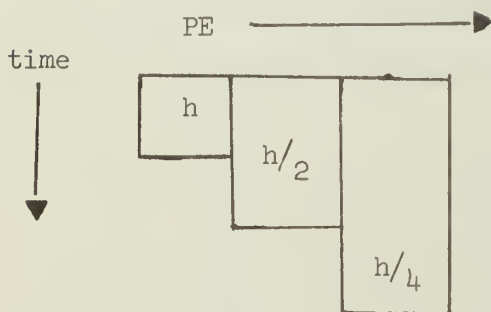
Each of these methods is stored as a skeleton ILLIAC IV program which is used to form the complete integrator output.

4. EVALUATION AND SUGGESTIONS FOR FURTHER DEVELOPMENT OF THE INTEGRATOR

As might well be expected, the efficiency of the ILLIAC IV program generated by the integrator system varies greatly on the type of input. One would expect it would not do as well on equations with highly complicated and diverse terms, and this proved to be true. Equations such as those involved in chemical kinetics problems where the terms differ only in two or three multiplies were handled very well, approaching the maximum efficiency for this algorithm.

The integrator is also not well suited to small systems of equations. There are many reasons for this. Conventional integration methods are rigidly sequential, so the work other than function evaluating can only use as many PE's as there are equations. Attempts to separate methods into usable parallel forms have been unsuccessful to date. This method thus leaves many of the PE's unoccupied when doing the method calculations for small systems. Also, with smaller systems, the ratio of these calculations to the function evaluation is high and so the overall integration becomes inefficient. One way of improving this might be to use some extrapolation method, such as Richardson's. A lower order, and hence faster method, could be used to do the integration, thus decreasing the integration-to-function evaluation ratio. Since this involves doing the integration at different stepsizes, more of the PE's could be performing useful work in parallel. However, due to the sequential nature of the methods, there would still be a lot

of PE's empty. For example, consider a three step extrapolation method halving the stepsize for each integration. The PE versus time graph would look as follows:



This represents the time involved in all the integration and function evaluation except for the evaluation of the extrapolation formula.

In this case 40 per cent of the machine ~~time~~ available is being used.

In the best case with only two stepsizes, 25 per cent of the machine time is unused, and the savings gained by an extrapolation from two cases would probably be negligible anyway. One is tempted to move part of the process time from the smaller stepsize to the empty spaces in the larger, but this cannot be done due to the aforementioned sequential nature of the methods. The starting values to begin the integration in the middle are not available until the integration has been performed up to that point. For these reasons extrapolation methods seem of doubtful value. There are other problems which affect small sets of equations. Since there are few terms, it can be expected that, on the average, there will only be a few that can be done in parallel. Thus the function evaluation may be inefficient as well.

With larger systems the results are significantly better. There are fewer wasted PE's during the integration operations. There are more terms to be compared and to fill up the machine. It cannot, however, be guaranteed that 100 per cent efficiency will be obtained even if the maximum number of terms can be done in parallel. This is due to the fact that if the number of terms is not an even multiple of 64, another term evaluation must be performed even if there is only one extra term. The following chart illustrates the maximum efficiency for different numbers of terms.

Maximum terms per PE	Worst case #terms	efficiency	Average #terms	efficiency
1	1	15%	32	50%
2	65	51%	96	75%
4	193	76%	224	87%
8	449	87%	480	94%
20	1217	95%	1248	98%

This clearly indicates some of the improvement with larger systems. Breaking up large terms into smaller ones might improve this although there will be some added overhead in keeping track of the extra terms and recombining to obtain the results.

Another item which causes inefficiency is the routing which must be done to recombine the partial sums after their evaluation. If the equations are roughly all the same length, the routing will be minimal. Strange distributions of long and short equations might greatly

increase the routing. An example might be a set where the first thirty equations have two terms each and the last thirty have twenty terms each. If these same equations were alternated - one short, one long, one short - the routing would be reasonable. It would be desirable to have the integrator system do this type of exchange before analyzing the equations. Of course, an intelligent user could do the same thing.

Another factor is the depth of analysis used to find terms that can be done in parallel. Highly complex terms are difficult to handle. Consider the following:

$$\begin{aligned} & Bx (Ax \sin (Y(1)) / 2 + Y (3) * 3) \\ & (Dx Y (2) - \sin (Y (6)) * Cx Y (2)) * 6 \end{aligned}$$

It is clear that many of these terms could be done in parallel, but highly sophisticated techniques would be needed to do this automatically. To do this high a level analysis on every term might take enormous amounts of computer time and might only be worthwhile on smaller systems.

As previously mentioned, another useful addition would be to have routines written so that a sine and a tangent, for example, could be done simultaneously. This would seem to be possible since the common methods for obtaining these functions are by using some approximating polynomial. It might be possible to use the same form polynomial with different constants to get different functions. It should be noted that sine and cosine are trivial to do in parallel since $\cos x = \sin (x + \pi/2)$.

It goes without saying that an integration method which is less sequential would be valuable in improving this system.

In summary, the integrator system works reasonably well on large systems of equations although there is still some improvement possible. For small systems it appears that a different algorithm would need to be used to get really good results.

APPENDIX A

SAMPLE INPUT PROBLEM FOR THE
INTEGRATOR-TRANSLATOR SYSTEM (OUTPUT LISTING)

TILIA IV TRANSLATOR WRITING SYSTEM
 GNDES COMPILER - VERSION 1 0 PROGRAM G /GNDES .
 APRIL 30, 1969 AT 2013 : 9.9

SLIST RSWD	*	200	1
# PATTERN 1	*	300	2
7(0) = A×Y(0)×Y(0)+R×Y(0)×Y(0)+C×Y(0)×Y(0) ;	*	400	3
7(1)=K1×Y(8)+A1×Y(1)+A6×Y(3)×Y(30)+A2×Y(1)+A3×Y(2)+A4×Y(7)+A5×Y(5)+000 ;	*	500	4
7(2)=B3×Y(1)×Y(9)+B2×Y(2)+B3×Y(2)+B4×Y(2)×Y(9)+05×Y(2)×Y(9)+K2×Y(5) ;	*	600	5
7(3)=C1×Y(3)×Y(9)+C2×Y(3)×Y(9)+C3×Y(3)+K3×Y(4)×Y(9)+K4×Y(1)+K4×Y(8)×Y(9)+ +K3×Y(8)×Y(3)+A1×Y(29)+K5×Y(1)×Y(2)+F6×Y(2)+K2×Y(3)×Y(5) ;	*	800	7
Z(4)=D1×Y(3)+D2×Y(4)×Y(9)+D3×Y(4)×Y(9) +K7×Y(8)×Y(4)+K8×Y(3) ;	*	900	8
7(5)=K3×Y(2)+E1×Y(2)×Y(9)+F2×Y(3)×Y(9)+F3×Y(4)×Y(9)+F4×Y(5)+F5×Y(5) ;	*	1000	9
7(6)=F1×Y(2)×Y(9)+F2×Y(3)×Y(9)+F3×Y(4)×Y(9)+F4×Y(6) +F3×Y(6)+ F1×Y(12)×Y(8)+A1×Y(17)×Y(26)+K1×Y(12)×Y(9)+K1×Y(9)+K4×Y(13)×Y(29) ;	*	1200	11
7(7)=G1×Y(6)+G2×Y(7)+G3×Y(7) ;	*	1300	12
7(8)=H1×Y(1) ;	*	1400	13
7(9)=I1×Y(10)+I2×Y(9)×Y(5)+I3×Y(9)×Y(7) ;	*	1500	14
7(10)=J1×Y(10)+J2×Y(10)×Y(7)+J3×Y(5)×Y(10)+K2×Y(8)+6.57849374 ;	*	1600	15
7(11)=K1×Y(10)+V3×Y(7)+K2×Y(9)×Y(7)+K3×Y(11) ;	*	1700	16
Z(12)=C1×Y(3)×Y(9)+C2×Y(3)×Y(9)+C3×Y(3)+A1×Y(9)+K4×Y(1)+K4×Y(8)×Y(9) ;	*	1800	17
7(13)=C1×Y(3)×Y(9)+C2×Y(3)×Y(9)+C3×Y(3)+ H3 ×Y(7)+K4×Y(1)+K4×Y(8)×Y(9) + +K5×Y(1)×Y(2)+K6×Y(2)+K2×Y(3)×Y(5)+A2×Y(16)×Y(23)+B5×Y(8)×Y(14) +C3×Y(26)×Y(15)+J1×Y(28)×Y(5)+A2×Y(22)+A4×Y(8)+K1×Y(8)×Y(2) ;	*	2100	20
7(14)=D1×Y(3)+D2×Y(4)×Y(9)+D3×Y(4)×Y(9) +K7×Y(8)×Y(4)+K8×Y(3) ;	*	2200	21
7(15)=K3×Y(2)+E1×Y(2)×Y(9)+F2×Y(3)×Y(9)+F3×Y(4)×Y(9)+E4×Y(5)+F5×Y(5) ;	*	2300	22
7(16)=F1×Y(2)×Y(9)+F2×Y(3)×Y(9)+F3×Y(4)×Y(9)+F4×Y(6) ;	*	2400	23
7(17)=G1×Y(6)+A1×Y(7)×Y(8)+G2×Y(7)+G3×Y(7) ;	*	2500	24
7(18)=H1×Y(1) ;	*	2600	25
7(19)=J1×Y(10)+A3×Y(1)×Y(23)+J2×Y(10)×Y(7)+J3×Y(5)×Y(10)+ 4.08-02×Y(3) ;	*	2700	26

$+ 5.344 \times Y(29) + Y(23) + 3.1416 \times Y(17) ;$	•	2800	27
$Z(20) = K1 \times Y(10) \times Y(7) + K2 \times Y(9) \times Y(7) + K3 \times Y(11) ;$	•	2900	28
$Z(21) = B3 \times Y(9) \times Y(10) + K2 \times Y(27) + J1 \times Y(10) + J2 \times Y(9) \times Y(5) + J3 \times Y(9) \times Y(7) ;$	•	3000	29
$Z(22) = C2 \times Y(3) \times Y(9) + C2 \times Y(3) \times Y(9) + C3 \times Y(3) \times Y(4) + K4 \times Y(3) \times Y(2) + K4 \times Y(8) \times Y(9) ;$	•	3100	30
$Z(23) = C2 \times Y(3) \times Y(9) + C2 \times Y(3) \times Y(4) + C3 \times Y(3) \times Y(8) + K4 \times Y(2) \times Y(25) + K4 \times Y(8) \times Y(9) +$ $+ F1 \times Y(1) \times Y(21) + K5 \times Y(2) \times Y(2) + K6 \times Y(2) + K2' \times Y(3) \times Y(5) + A1 \times Y(8) + F4 \times Y(9) \times Y(14) ;$	•	3200	31
$+ C1 \times Y(8) \times Y(9) + C2 \times Y(3) + C3 \times Y(2) + C6 \times Y(9) \times Y(9) ;$	•	3300	32
$Z(24) = D2 \times Y(3) + C2 \times Y(4) \times Y(9) + C3 \times Y(4) \times Y(4) + K7 \times Y(8) \times Y(4) + K8 \times Y(3) ;$	•	3400	33
$Z(25) = K3 \times Y(2) + F2 \times Y(2) \times Y(9) + F2 \times Y(3) \times Y(4) + F3 \times Y(4) \times Y(4) + F4 \times Y(5) + F5 \times Y(5) ;$	•	3500	34
$Z(26) = F2 \times Y(2) \times Y(9) + F2 \times Y(3) \times Y(4) + F3 \times Y(4) \times Y(9) + F4 \times Y(6) ;$	•	3600	35
$Z(27) = G2 \times Y(6) + G2 \times Y(7) + G3 \times Y(7) ;$	•	3700	36
$Z(28) = H2 \times Y(2) ;$	•	3800	37
$Z(29) = J2 \times Y(2) + J2 \times Y(20) \times Y(7) + J3 \times Y(5) \times Y(20) + A2 \times Y(8) + J1 \times Y(1) \times Y(9) + 4.0 ;$	•	3900	38
$Z(30) = J2 \times Y(20) + J2 \times Y(9) \times Y(5) + J3 \times Y(4) \times Y(7) ;$	•	4000	39
$Z(1) = 118.0 ;$	•	4100	40
$Z(2) = 78.0 ;$	•	4200	41
$Z(3) = 37.0 ;$	•	4300	42
$Z(4) = 14900.0 ;$	•	4400	43
$Z(5) = 57.0 ;$	•	4500	44
$Z(6) = 29.0 ;$	•	4600	45
$Z(7) = 430.0 ;$	•	4700	46
$Z(8) = 0.0 ;$	•	4800	47
$Z(9) = 1.0 ;$	•	4900	48
$Z(10) = 4.0 ;$	•	5000	49
$Z(11) = 0.0 ;$	•	5100	50
$Z(12) = 3.5645 ;$	•	5200	51
$Z(13) = 34567.89 ;$	•	5300	52
$Z(14) = 0 ;$	•	5400	53
$Z(15) = 1 ;$	•	5500	54
$Z(16) = .0000005 ;$	•	5600	55
$Z(17) = 3.00987654 ;$	•	5700	56
	•	5800	57

$Z_0(18) = .903$;
 $Z_0(19) = 4.675$;
 $Z_0(20) = 0$;
 $Z_0(21) = 345.67452-1$;
 $Z_0(22) = 3356345.0$;
 $Z_0(23) = 3.0$;
 $Z_0(24) = 0$;
 $Z_0(25) = 0$;
 $Z_0(26) = 0$;
 $Z_0(27) = 0$;
 $Z_0(28) = 0$;
 $Z_0(29) = 0$;
 $Z_0(30) = 0$;
 $A_1 = -1.0$;
 $A_2 = -2.0$;
 $A_3 = .05$;
 $A_4 = 0.1$;
 $A_5 = 0.56$;
 $A_6 = 3.1415$;
 $R_1 = 1.0$;
 $R_2 = -0.05$;
 $R_3 = -1.5$;
 $R_4 = -0.03$;
 $R_5 = -0.03$;
 $C_1 = -0.02$;
 $C_2 = -0.02$;
 $C_3 = -3.0$;
 $D_1 = 3.0$;
 $D_2 = -0.005$;
 $D_3 = -0.005$;
 $F_1 = 0.01$;

*	5900	58
*	6000	59
*	6100	60
*	6200	61
*	6300	62
*	6400	63
*	6500	64
*	6600	65
*	6700	66
*	6800	67
*	6900	68
*	7000	69
*	7100	70
*	7200	71
*	7300	72
*	7400	73
*	7500	74
*	7600	75
*	7700	76
*	7800	77
*	7900	78
*	8000	79
*	8100	80
*	8200	81
*	8300	82
*	8400	83
*	8500	84
*	8600	85
*	8700	86
*	8800	87
*	8900	88

F2= 0.02 ;	• 9000	89
F3= 0.005 ;	• 9100	90
F4= -0.54 ;	• 9200	91
F5= -0.033 ;	• 9300	92
F1= 0.03 ;	• 9400	93
F2= 0.02 ;	• 9500	94
F3= 0.005 ;	• 9600	95
F4= -2.0 ;	• 9700	96
G1= 2.0 ;	• 9800	97
G2=-0.01 ;	• 9900	98
G3=-0.033 ;	• 10000	99
H1= 2.0 ;	• 10100	100
H2= 5.084978775 ;	• 10200	101
I2=-0.01000 ;	• 10300	102
I3=-0.00733 ;	• 10400	103
J1=-0.50 ;	• 10500	104
J2=-0.00758 ;	• 10600	105
J3=-0.00750 ;	• 10700	106
K1= 0.00758 ;	• 10800	107
K2= 0.00733 ;	• 10900	108
K3=-1.0 ;	• 11000	109
K4=-3.58497364 ;	• 11100	110
K5= -5456.48586 ;	• 11200	111
K6= -345.56778=0.1 ;	• 11300	112
K7= 45677.999 ;	• 11400	113
K8= 45.566859 ;	• 11500	114
K9= 4566.4907 ;	• 11600	115
T0 = 0 ;	• 11700	116
TF =1 ;	• 11800	117
TP =0.0,0.25,0.5,0.75,1.0 ;	• 11900	118
STEP = 0.05 ;	• 12000	119

APPENDIX B

ILLIAC IV ASSEMBLY LANGUAGE OUTPUT
PRODUCED BY THE INTEGRATOR FOR THIS SAMPLE PROBLEM

GODES /RIGASK LISTED BY SNAP ON APRIL 30, 1969 AT 20:39:44.8

[illegible]

FILL	128;	00 05600	50
DATA		00 05400	59
	-1.0,0.05,600,-1.5,0.00233,-3.0,-3.58697364,	00 06000	60
	-5456.48586,3.0,45677.999,0.03,-0.56,0.02,0,005,-1,0,	00 06100	61
	2.0,2.0,-0.00233,-0.00250,0.00058,-1.0,-3.0,	00 06200	62
	-3.58697364,-3.0,-3.58697364,0.00233,-3.0,0.1,-0.005,	00 06300	63
	45.566859,0.02,-0.03,0.005,-1.0,2.0,-0.00058,5.346,	00 06400	64
	0.00233,0.00233,-0.00233,-3.0,-0.02,-3.58697364,	00 06500	65
	-5456.48586,-1.0,-0.005,-0.005,45677.999,0.02,-0.56,	00 06600	66
	0.02,-0.01,5.566973775,-0.00250,4.0,-0.00233,0.0,0.0,0,	00 06700	67
	0;	00 06800	68
FILL	128;	00 06900	69
DATA		00 07000	70
	0.00058,-2.0,0.56,-0.05,-0.03,-0.02,-3.58697364,-1.0,	00 07100	71
	0.00233,-0.005,-1.0,0.005,0.03,-2.0,-1.0,-3.58697364,	00 07200	72
	-0.033,-0.01000,-0.00058,6.57549274,0.00233,-0.02,	00 07300	73
	-3.58697364,-0.02,-3.58697364,-345.56778,-04,-0.03,	00 07400	74
	-0.03,0.045,77.999,0.03,-0.56,0.02,0.033,0.05,	00 07500	75
	4.08,-0.2,-0.00058,-1.5,-0.01000,-0.02,-3.58697364,-3.0,	00 07600	76
	0.03,-0.00233,-0.02,-0.005,-1.0,0.005,0.02,-2.0,	00 07700	77
	-0.033,-0.00058,-0.56,-0.01000,0.0,0.0,0,0;	00 07800	78
FILL	128;	00 07900	79
DATA		00 08000	80
	0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,	00 08100	81
	0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,	00 08200	82
	0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0,0;	00 08300	83
FILL	128;	00 08400	84
DATA		00 08500	85
		00 08600	86
	3, 1, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11, 12, 13, 14,	00 08700	87
	15, 16, 17, 18, 19, 20, 21, 22, 23, 24, 25, 26, 27, 28, 29,	00 08800	88
	30, 31, 32, 33, 34, 35, 36, 37, 38, 39, 40, 41, 42, 43,	00 08900	89
	44, 45, 46, 47, 48, 49, 50, 51, 52, 53, 54, 55, 56, 57,	00 09000	90
	58, 59, 60, 61, 62, 63, 64, 65, 66, 67, 68, 69, 70,	00 09100	91
	71, 72;	00 09200	92
FILL	128;	00 09300	93
DATA		00 09400	94
	0, 7, 9, 2, 9, 7, 3, 2, 9, 3, 9, 5,	00 09500	95
	0, 8, 9, 7, 10, 10, 8, 7, 9, 9, 9, 9,	00 09600	96
	2, 23, 5, 2, 2, 9, 9, 6, 7, 10, 10,	00 09700	97
	31, 11, 10, 4, 2, 9, 9, 2, 14, 2, 9, 3,	00 09800	98
	9, 5, 9, 7, 20, 8, 20, 0, 0, 0, 0, 0,	00 09900	99
	0, 1;	00 10000	100
FILL	128;	00 10100	101
DATA		00 10200	102
	0, -1, 0, -1, 0, -1, -2, -3, -2, -3, -2, -3,	00 10300	103
	-2, -3, -4, -3, 0, 1, 0, 1, 2, 1, 2, 1,	00 10400	104
	0, -1, -2, -3, -2, -1, -2, -1, -2, -1, 2, 1,	00 10500	105
	0, 1, 2, 3, 2, 3, 2, 1, 0, -1, 0, -1,	00 10600	106
	0, -1, 0, 1, 4, 3, 2, 3, 0, 0, 0, 0,	00 10700	107
	0, 0;	00 10800	108
FILL	128;	00 10900	109
DATA		00 11000	110
	0, 0, 0, 0, 0, 0, 8, 5, 0, 8, 2, 0,	00 11100	111
	3, 0, 12, 0, 0, 9, 5, 0, 0, 8, 0,	00 11200	112
	4, 3, 20, 0, 4, 0, 3, 0, 4, 7, 0, 10,	00 11300	113
	27, 9, 0, 9, 3, 3, 2, 2, 0, 0, 0, 8,	00 11400	114
	2, 0, 3, 0, 0, 5, 0, 5, 0, 0, 0, 0,	00 11500	115
	0, 0;	00 11600	116
FILL	128;	00 11700	117
DATA		00 11800	118

TAGS:

LDA	XC01X	00	18000	180
STA	XS1X	00	18100	181
LDA	XT1X	00	18200	182
STA	XT-11X	00	18300	183
STA	XK1X	00	18400	184
STA	XKTFM1X	00	18500	185
*.....END OF DATA LOADING.....				
LIT(3)	=COUNTS	00	18600	186
STL(3)	=COUNTS	00	18700	187
* THIS NEXT SECTION IS USED IF A DELAY IS USED. (X(T-TAU))				
* THE FOLLOWING PART OF THE PROGRAM IS THE RUNGEKUTTA CALCULATIONS				
NEWT:	STL(2)	=TEMP21X	00	18800
	LIT(1)	=IAT1X	00	18900
	STL(3)	=TEMP31X	00	19000
	SILT(2)	=EVAL1X	00	19100
	FXCHL(2)	=ICR1X	00	19200
	SKIP	=01X	00	19300
	LDA	DELT1X	00	19400
	MIRN	=C11X	00	19500
	ADRN	=I01X	00	19600
	STA	=I01X	00	19700
	LDA	SUM1X	00	19800
	STA	SUM1X	00	19900
	MIRN	=C11X	00	20000
	MIRN	DELT1X	00	20100
	ADRN	XK1X	00	20200
	STA	XKTFM1X	00	20300
	SILT(2)	=EVAL1X	00	20400
	FXCHL(2)	=ICR1X	00	20500
	SKIP	=01X	00	20600
	LDA	SUM1X	00	20700
	STA	SUM21X	00	20800
	MIRN	DELT1X	00	20900
	MIRN	=C11X	00	21000
	ADRN	XK1X	00	21100
	STA	XKTFM1X	00	21200
	SILT(2)	=EVAL1X	00	21300
	FXCHL(2)	=ICR1X	00	21400
	SKIP	=01X	00	21500
	LDA	SUM1X	00	21600
	STA	SUM31X	00	21700
	MIRN	DELT1X	00	21800
	ADRN	XK1X	00	21900
	STA	XKTFM1X	00	22000
	LDA	DELT1X	00	22100
	MIRN	=C11X	00	22200
	ADRN	=I01X	00	22300
	STA	=I01X	00	22400
	SILT(2)	=EVAL1X	00	22500
	FXCHL(2)	=ICR1X	00	22600
	SKIP	=01X	00	22700
	LDA	SUM21X	00	22800
	ADRN	SUM31X	00	22900
	LIT(2)	=IWN1X	00	23000
	MIRN	=C21X	00	23100
	ADRN	SUM1X	00	23200
	ADRN	SUM1X	00	23300
	MIRN	DELT1X	00	23400
	LIT(1)	=ICR1X	00	23500
	MIRN	=C11X	00	23600
	ADRN	XK1X	00	23700
			00	23800
			00	23900
			00	24000

STA	X41X		00	24100	241	
STA	X4FM1X		00	24200	242	
L0L(2)	.TFMP21X		00	24300	243	
TYLTM(2)	.NFWT1X	IF NOT TO BE SAVED STEP	00	24400	244	
STA	X1(3)1X		00	24500	245	
L0L(2)	.INC1X		00	24600	246	
L0L(3)	.TFMP31X		00	24700	247	
TYLTM(3)	.NFWT1X	LOOP UNTIL END	00	24800	248	
JUMP	U01X	*****	00	24900	249	
SKIP	.01X		00	25000	250	
* END OF RUNGE KUTTA CALCULATIONS ROUTINE			00	25100	251	
*			00	25200	252	
* BEGINNING OF FUNCTION EVALUATION SUBROUTINE.			00	25300	253	
FVAL:	EXCHL(2)	.S021X	SAVE CARS	00	25400	254
	EXCHL(3)	.S031X		00	25500	255
	STL(0)	.S0111X		00	25600	256
	STL(1)	.S0121X		00	25700	257
	SLIT(2)	.R0TF1X	CALL NOTE	00	25800	258
	EXCHL(2)	.SICR1X		00	25900	259
	SKIP	.01X		00	26000	260
	LDA	.01X		00	26100	261
	STA	SUM1X	ZERO SUM	00	26200	262
FFIJ:	STL(2)	.S071X	*****	00	26300	263
	L0L(2)	.ROUND1X	RANGE OF ROUTES LOADED	00	26400	264
	LDA	.01X	ZERO OUT TSUM LOCATIONS	00	26500	265
INIT:	CSUR(2)	.LEFT1X		00	26600	266
	STA	TSUM(2)1X		00	26700	267
	CAUD(2)	.LEFT1X		00	26800	268
	TYLTM(2)	.INIT1X		00	26900	269
	LDA	CONST+ 0;		00	27000	270
	LDX	TAGS+ 0;		00	27100	271
	MLRN	*XS ;		00	27200	272
	LDX	TAGS+ 1;		00	27300	273
	MLRN	*XS ;		00	27400	274
	LDX	TAGS+ 2;		00	27500	275
	ADRN	*TSUM;		00	27600	276
	STA	*TSUM;		00	27700	277
	LDA	CONST+ 1;		00	27800	278
	LDX	TAGS+ 3;		00	27900	279
	MLRN	*XS ;		00	28000	280
	LDX	TAGS+ 4;		00	28100	281
	MLRN	*XS ;		00	28200	282
	LDX	TAGS+ 5;		00	28300	283
	ADRN	*TSUM;		00	28400	284
	STA	*TSUM;		00	28500	285
	LDA	CONST+ 2;		00	28600	286
	LDX	TAGS+ 6;		00	28700	287
	MLRN	*XS ;		00	28800	288
	LDX	TAGS+ 7;		00	28900	289
	MLRN	*XS ;		00	29000	290
	LDX	TAGS+ 8;		00	29100	291
	ADRN	*TSUM;		00	29200	292
	STA	*TSUM;		00	29300	293
	LDA	.01X	ZERO A REGISTERS	00	29400	294
	L0L(3)	.ROUND1X	ROUTING ROUND	00	29500	295
FINS:	L0L(2)	.S031X		00	29600	296
	CSUR(2)	.LEFT1X	SUR LEFT MAX TO GET NEGATIVE ROUTES	00	29700	297
	LDR	TSUM(2)1X	RGR= PARTIAL SUMS	00	29800	298
	CAND(2)	.V631X	MUST ON ROUTE OF 43 INSTEAD OF -1	00	29900	299
	RTL	0(2)1X	ROUTE PARTIAL SUMS TO CORRECT PF	00	30000	300
	ADRN	.3R1X	ADD TO A REGISTERS	00	30100	301

APPENDIX C

BNF INPUT SPECIFICATIONS FOR THE INTEGRATOR

(INPUT TO THE TWS SYSTEM ALONG WITH SEMANTICS
TO PRODUCE THE INTEGRATOR-TRANSLATOR. THE
SEMANTICS ARE TOO LONG TO BE LISTED HERE)

GNDFS /SYNTAX LISTED BY SNAP ON APRIL 30, 1969 AT 20139143.3

/GNDFS .	00 00100	1
PRINT ALLTAB;	00 00200	2
FPARSER)	00 00300	3
#Y,T,Z,LIST,TF,70,FND,LOG,T0,TP,CUS,EXP,VARSTFP,HUNGF,SIN,EXTRAP;	00 00400	4
<PROG>:= <STATE> #5 / <TERM> <SFP> <STATE> #5;	00 00500	5
<SFP> := #1 / <>;	00 00600	6
<STATE> := <DIFFQN> #10 / <INIT> #15 / <TIMIT> #20 / <TFINAL> #25 /	00 00700	7
<TPRINT> #30 / <CONSTDEF> #35 / <CONTRUL> #40 /	00 00800	8
#STEP = <REALNUM> #41 / #FND #1 / #FRROR = <REALNUM> #61 ;	00 00900	9
<CONTRUL> := # <CONWRD> #44 / <CONTRUL> <CONWRD> #43 ;	00 01000	10
<CONWRD> := #DEFBIG #42 / #TEST #39 / #CONRRWF #38 / #RUNGF #37 /	00 01100	11
#EXINSERT = <*N> #4 / / #CONIEINSERT = <*N> #45 / #FINMONE #57	00 01200	12
/ #FRRURCHECK #54 / #CONSTRIP #52 / #PATTERN #60;	00 01300	13
<DIFFQN>:= #Z(<INDEX>) #7 = <FXW>;	00 01400	14
<INDEX>:= <*N> #36 ;	00 01500	15
<EXP> := <TERM> #34 / <ANP> <TERM> #33 / <EXP> <ANP> <TERM> #32 ;	00 01600	16
<TERM>:= <FACTOR> #31 / <TERM> <MNP> <FACTOR> #20;	00 01700	17
<ANP> := + #28 / - #27 ;	00 01800	18
<FACTOR> := <PRIMARY> #26 / <FACTOR> * #48 <PRIMARY> #24 ;	00 01900	19
<MNP> := x #23 / #/ #22 ;	00 02000	20
<PRIMARY>:= <REALNUM> #21 / <*I> #19 / #Y(<INDEX>) #18 / #T #17 / <FUNCT> #16	00 02100	21
/ (#2 <FXW>) #14 ;	00 02200	22
<REALNUM>:= <*N> #40 <FRACTION> <EXPONENT> #54 / . <*N> #56	00 02300	23
<EXPONENT> #54;	00 02400	24
<FRACTION>:= . <*N> #50 / . / <> ;	00 02500	25
<EXPONENT>:= # <EXPSTGN> <*N> #51 / <>;	00 02600	26
<STGN>:= + / - #55 / <>;	00 02700	27
<EXPSTGN>:= + / - #52 / <>;	00 02800	28
<FUNCT> := <FUNID> (<FXW>) #13 ;	00 02900	29
<FUNID> := #SIN #12 / #COS #11 / #LOG #9 / #EXP #8 ;	00 03000	30
<INIT>:= #70(<INDEX>) = <STGN> <REALNUM>;	00 03100	31
<TIMIT>:= #T0 = <STGN> <REALNUM>;	00 03200	32
<TPRINT>:= #TP = <STGN> <REALNUM> #46 / <TPRINT> <STGN> <REALNUM> #7 /	00 03300	33
#TP = (<REALNUM>) #6 ;	00 03400	34
<TFINAL>:= #TF = <STGN> <REALNUM>;	00 03500	35
<CONSTDEF>:= <*T> #4 = <STGN> <REALNUM>;	00 03600	36
	00 03700	37

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13. ABSTRACT This report describes the design and implementation of an automatic integrator for systems of ordinary differential equations for use on the ILLIAC IV computer system. This integrator accepts a simple statement of the differential equations and their initial values and produces an ILLIAC IV Assembly Language code which can be used to solve them. The intent is to produce a code which will be efficient on a large parallel computer.			

14. KEY WORDS		LINK A		LINK B		LINK C	
		ROLE	WT	ROLE	WT	ROLE	WT
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